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Introduction

 Demand for cloud-based processing frameworks continues to grow Cloud applications is becoming increasingly mission-critical and deadline sensitive, especially in shared clusters Cloud providers seek efficient techniques to meet the SLA • A few slow tasks, called stragglers, could significantly impact job execution time Launching extra attempts (clones) for each task upon submission can mitigate stragglers

 This work proposes Shed, an optimization framework that leverages dynamic cloning to jointly maximize jobs' Probability of Completion before Deadline (PoCD) by fully utilizing the available resources

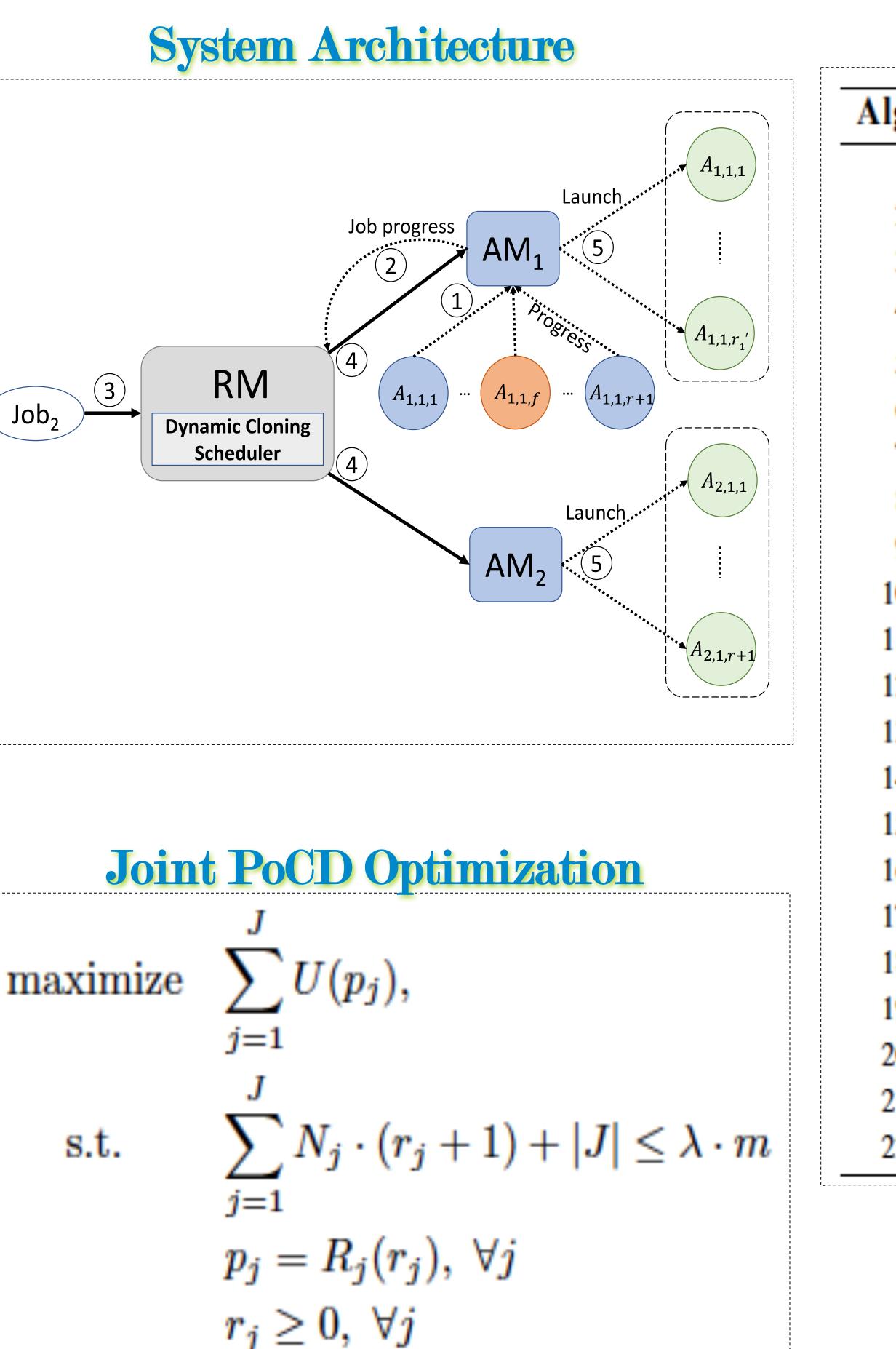
System Model

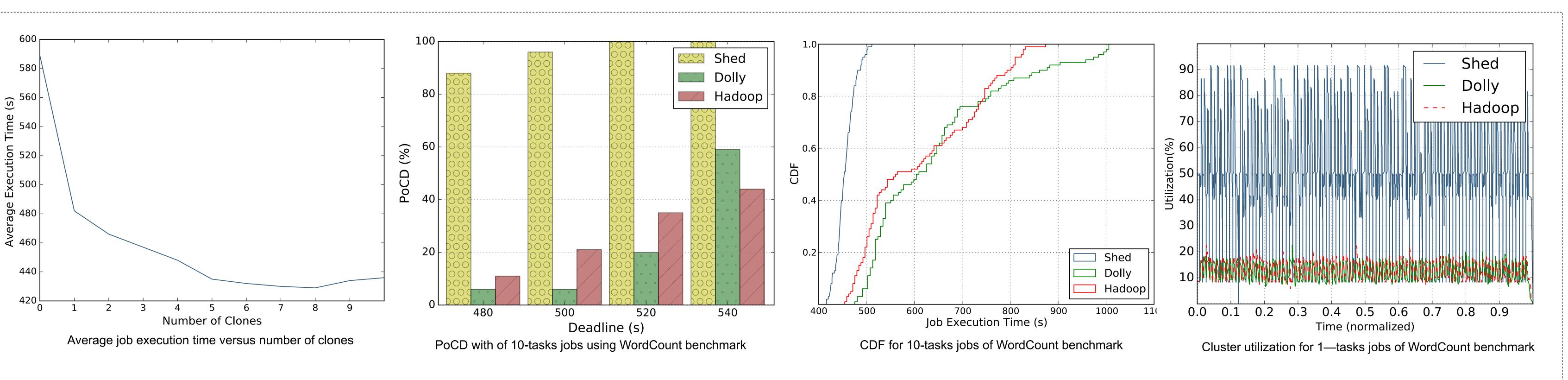
jobs submitted to the Consider Mapreduce processing framework. Each job j is associated with a deadline Dj that is determined by a user. Each job j consists of Nj tasks, and it is considered successful if all its Nj tasks are executed and completed before the job deadline Dj. Let Tj denote job j's completion time, and Tj,i for i = 1, ..., Nj be the (random) completion times of tasks belongs to job j. Any task whose execution time exceeds the deadline is considered a straggler. Our dynamic cloning approach mitigate the effect of stragglers by proactively launching rj extra attempts for each task. A task is finished once any one of the rj + attmepts finishes execution, and then the other copies are killed. Thus, taks i's completion time Tj,i is determined by the completion time of the fastes attempt, i.e.,

 $T_{j,i} = \min_{k=1,...,r_j+1} T_{j,i,k}, \ \forall i, j.$

Shed: Optimal Dynamic Cloning to Meet **Application Deadlines in Cloud**

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Proposed Online Algorithm

Igorithm 1: Proposed Online Algorithm
1: Upon submission of a new job:
2: Kill all jobs which missed their deadlines
3: $J = \{j_1, j_2, j_3, \ldots\}$
4: if $ J == 1$ then
5: $r_{\max} = \left\lfloor \frac{\lambda \cdot m - N_1 - 1}{N_1} \right\rfloor$ 6: $r_1 = r_{\max}$
6: $r_1 = r_{\text{max}}$
7: else
8: $r_j = 0 \forall j$
9: $\omega = 0$
10: $\kappa = \lambda \cdot m - \sum_{j=1}^{J} N_j - J $
11: Calculate $R_j \forall j$
12: while $J \neq \{\emptyset\}$ do
13: $j' = \arg\min_j \{R_j\}$
14: if $N_{j'} + \omega > \kappa$ then
15: $J = J - \{j'\}$
16: else
17: $r_{j'} = r_{j'} + 1$
18: $\omega = \omega + N_{j'}$
19: Calculate $R_j \forall j$
20: end if
21: end while
22: end if

Shed, an In this work, we propose that optimization framework leverages dynamic cloning to jointly maximize PoCD and cluster utilization. We also present an online scheduler that dynamically optimize resources upon new job arrival. Our solution includes an online greedy algorithm to find the optimal number of clones needed for each job. Our results show, in some cases, can achieve 100% PoCD that Shed to Dolly and Hadoop with compared speculation. The proposed algorithm is able to achieve more than 90% utilization of available cloud resources, whereas Dolly and Hadoop achieves only about 22%.

Evaluation

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Theorems

$R_{\rm su} = \left[1 - \left(\frac{t_{\rm min}}{D_i}\right)^{\beta \cdot (r_j + 1)} \right]^{N_j}$ $R_{\rm ru} = \left[1 - \left(\frac{(1 - \phi_i) t_{\rm min}}{D_j - \tau_j} \right)^{\beta \cdot (r_j + 1)} \right]^{N_j}$

Conclusion



• J. Dean *et al*, 2008

• G. Ananthanarayanan *et al*, 2013